

## 4. MIXED STRATEGY AND DOMINANCE

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### Exercise 1.

#### ■ Nash Equilibrium

		Player 2	
		<i>E</i> [ <i>z</i> ]	<i>F</i> [ <i>1 - z</i> ]
Player 1	<i>A</i> [ <i>x</i> ]	5 4	5 4
	<i>B</i> [ <i>y</i> ]	8 3	1 9
	<i>C</i> [ <i>1 - x - y</i> ]	3 6	7 2

Matrix A-1

As you can see from Matrix A-1, there is no pure Nash equilibrium. To find the mixed Nash equilibria, we suppose Player 1 plays

$$\left\{ \begin{array}{l} A \text{ with probability } x, \\ B \text{ with probability } y, \\ C \text{ with probability } 1 - x - y, \end{array} \right.$$

that is, we suppose that Player 1's mixed strategy is

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Similarly, we can assume that Player 2 takes the following mixed strategy [according to Matrix A-1]:

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**Step 1** If Player 2 plays a mixed strategy, then he must be indifferent between  $E$  or  $F$ ; hence Player 2 must get the same payoffs from playing  $E$  and  $F$ . His payoff from playing  $E$  is

$$\underline{\hspace{10em}};$$

his payoff from playing  $F$  is

$$\underline{\hspace{10em}}$$

These two payoffs are equal, so we know  $x = 1 - 5y/2$ .

**Step 2** According to the above equation, there are three cases:

- $x = 0$ . In this case, Player 1 does not play  $A$ .
- $x \in (0, 1)$ , which means that  $y \in (0, 2/5)$ . In this case, Player 1 plays  $A$  and  $B$  randomly [he may or may not play  $C$ ; we have no idea at this stage].
- $x = 1$ , which means that  $y = 0$ . In this case, Player 1 plays  $A$  deterministically.

**Step 3** If  $x = 1$ , then  $A$  is Player 1's best response against to Player 2's mixed strategy  $(z, 1 - z)$ ; thus,

$$\begin{cases} A \text{ is better than } B \Rightarrow \underline{\hspace{10em}} \\ A \text{ is better than } C \Rightarrow \underline{\hspace{10em}} \end{cases}$$

i.e.,  $1/2 \leq z \leq 4/7$ .

**Step 4** Check by yourself that

$$\left\{ (A, z \cdot E + (1 - z) \cdot F) \mid \frac{1}{2} \leq z \leq \frac{4}{7} \right\}$$

is a set of Nash equilibria.

**Step 5** If  $x = 0$ , then  $y = 2/5$  and  $1 - x - y = 1 - 0 - 2/5 = 3/5$ ; that is, Player 1 plays  $B$  with probability  $2/5$ , and plays  $C$  with probability  $3/5$ . In this case, Player 1 must get the same payoffs from playing  $B$  and  $C$ , i.e.,

$$8 \cdot z + 1 \cdot (1 - z) = 3 \cdot z + 7 \cdot (1 - z),$$

which solves for  $z = 6/11$ .

While getting

$$8 \times \frac{6}{11} + 1 \times \left(1 - \frac{6}{11}\right) = \frac{53}{11}$$

from playing  $B$  or  $C$  randomly, Player 1 can guarantee 5 if he plays  $A$  deterministically. Since  $5 > 53/11$ , we know that  $x = 0$  is impossible.

**Step 6** We now check the last possibility:  $0 < x < 1$ . In this case, Player 1 gets the same payoffs from  $A$  and  $B$  [notice that in this case both  $x$  and  $y$  are greater than 0]. Hence,

$$5 \cdot z + 5 \cdot (1 - z) = 8 \cdot z + 1 \cdot (1 - z),$$

which solves for  $z = 4/7$ .

If Player 2 adopts the mixed strategy

$$\left(\frac{4}{7}E + \frac{3}{7}F\right),$$

Player 1's expected payoff from  $C$  is

$$3 \times \frac{4}{7} + 7 \times \frac{3}{7} = \frac{33}{7} < 5,$$

which means that Player 1 will not use  $C$ , or equivalently,

$$1 - x - y = 0.$$

Therefore,  $x + y = 1$ . The proceeding equation with  $x = 1 - \frac{5}{2}y$  yield

$$x = 1 \quad \text{and} \quad y = 0.$$

A contradiction [remember that in this step we assume  $0 < x < 1$  and  $0 < y < 1$ ].

**Step 7** From Step 1–Step 6, we conclude that the set of Nash equilibria in this game is

#### ■ Undominated Strategy

All Nash equilibria are undominated strategies.

**Exercise 2.** See the solution.

**Exercise 3.**

#### ■ The Game in Matrix C-1

**Step 1** For Player 1,  $B$  is strictly dominated. We delete  $B$  and this produce the following Matrix C-3.

		Player 2	
		L	R
Player 1	T	2	2
	C	2	$\frac{1}{2}$
	B	0	1

2	2
3	0
0	1

		Player 2	
		L	R
Player 1	T	2	2
	C	3	$\frac{1}{2}$
	B	0	1

2	2
3	0
0	1

Matrix C-1

Matrix C-2

		Player 2	
		L	R
Player 1	T	2	2
	C	2	$\frac{1}{2}$
	B	0	1

Matrix C-3

**Step 2** In the game Matrix C-3, we can find the pure Nash equilibria easily. Now let Player 1's mixed strategy be  $x \cdot T + (1 - x) \cdot C$ , and Player 2's mixed strategy be  $y \cdot L + (1 - y) \cdot R$ .

**Step 3** As usually, Player 1 is indifferent between  $T$  or  $C$ , so

$$2 = 2y + \frac{1}{2}(1 - y) \Rightarrow y = 1,$$

which means that if Player 2 plays  $L$  with certainty, then Player 1 is indifferent between  $T$  or  $C$  [you can see this from the Matrix C-3]. Hence, given Player 2's strategy  $L$ , Player 1's best response is

$$x \cdot T + (1 - x)C, \quad x \in [0, 1].$$

**Step 4** We also need to check that given Player 1's mixed strategy  $x \cdot T + (1 - x)C$ ,  $x \in [0, 1]$ , Player 2's best response is  $L$ . This is true because

$$2x + 3(1 - x) = 3 - x \geq 2x.$$

**Step 5** With the same logic, you can find the following set of Nash equilibria:

$$\left\{ (T, y \cdot L + (1 - y)R) \mid y \in [0, 1] \right\}.$$

■ The Game in Matrix C-2

**Step 1** Delete the strictly dominated strategy  $B$ , and we get Matrix C-4.

		Player 2	
		$L$ [ $y$ ]	$R$ [ $1 - y$ ]
Player 1	$T$ [ $x$ ]	2	2
	$C$ [ $1 - x$ ]	3	$\frac{1}{2}$ 0

Matrix C-4

**Step 2** The pure Nash equilibria are  $(T, R)$  and  $(C, L)$ .

**Step 3** Consider the mixed strategies [sometimes I do not write the mixed strategies explicitly for simplicity. You can find them from the corresponding matrix]. There is a simple way to find the mixed Nash equilibria. For Player 2, he would like to play a mixed strategy  $(y \cdot L + (1 - y) \cdot R)$  if and only if  $x = 1$ .

**Step 4** For Player 1, he would like to play  $T$  with certainty if and only if

$$2 \geq 3y + \frac{1}{2}(1 - y) \Rightarrow y \leq \frac{3}{5}.$$

**Step 5** We need not to consider any of Player 1's mixed strategy since if  $x < 1$ , Player 1's best response is  $L$ , and if Player 2's strategy is  $L$ , Player 1's best response is  $C$ .

**Step 6** We thus know that the set of Nash equilibria is

$$\left\{ (T, y \cdot L + (1 - y) \cdot R) \mid y \in \left[ 0, \frac{3}{5} \right] \right\} \text{ and } \{(C, L)\}.$$

**Exercise 4.**

**Step 1** If Player 1 uses the following mixed strategy

$$x \cdot A + y \cdot B + (1 - x - y) \cdot C,$$

where  $x > 0$ , and  $1 - x - y > 0$ . We first suppose that  $x \leq 1 - x - y$ , then we can rewrite his mixed strategy as [Consider the case of  $x \geq 1 - x - y$  by yourself]

$$2x \cdot \left( \frac{1}{2}A + \frac{1}{2}C \right) + y \cdot B + (1 - 2x - y) \cdot C.$$

For simplicity, we denote the above mixed strategy as  $CM$ .

		Player 2		
		<i>D</i>	<i>E</i>	<i>F</i>
Player 1	<i>A</i> [ <i>x</i> ]	1 2	3 0	0 3
	<i>B</i> [ <i>y</i> ]	1 1	2 2	2 0
	<i>C</i> [ $1 - x - y$ ]	1 2	0 3	3 0

**Step 2** Because  $\frac{1}{2}A + \frac{1}{2}C$  is dominated by *B*, we know immediately that *CM* is dominated by

$$(2x + y) \cdot B + (1 - 2x - y) \cdot C.$$

**Step 3** Step 1 and Step 2 thus show that any of Player 1's mixed strategies putting positive probabilities on *A* and *C* are dominated.

**Step 4** There are two cases under which Player 1's mixed strategy

$$x \cdot A + y \cdot B + (1 - x - y) \cdot C$$

cannot be rewritten as *CM*:  $x = 0$  or  $1 - x - y = 0$ . We are not interested in the former case [ $x = 0$ ] because it is not part of a Nash equilibrium [remember that  $x \geq \frac{1}{3}$  in any Nash equilibria].

**Step 5** Thus we need only consider the case of  $1 - x - y = 0$ , which means that Player 1's mixed strategies are

$$x \cdot A + (1 - x) \cdot B, \quad x \geq \frac{1}{3}.$$

**Step 6** Note that  $1 - x - y = 0$  implies that

$$y = 1 - x.$$

Combining this equation with the following equation

$$y \leq 2 - 3x,$$

we have

$$1 - x \leq 2 - 3x,$$

which solves for

$$x \leq \frac{1}{2}.$$

**Step 7** We have proved through Step 1—Step 6 that for Player 1, if his strategy is undominated and is part of Nash equilibrium, then his strategy must be

$$x \cdot A + (1 - x) \cdot B, \quad \frac{1}{3} \leq x \leq \frac{1}{2}.$$

Our final step is to show that the above strategies are really undominated.

**Step 8** Suppose that there exist a strategy

$$a \cdot A + b \cdot B + (1 - a - b) \cdot C$$

which dominates  $x \cdot A + (1 - x) \cdot B$ , where  $a, b$  are probabilities satisfying  $a + b \leq 1$ . We need to solve the following system of inequalities:

$$\left\{ \begin{array}{ll} 3a + 2b \geq 2 + x, & \text{if } s_2 = E \\ 3 - 3a - b \geq 2 - 2x, & \text{if } s_2 = F \\ \text{At least one of the above two inequalities holds strictly,} \\ a + b \leq 1, \\ 0 \leq a \leq 1, \\ 0 \leq b \leq 1. \end{array} \right.$$

There is no solution for the above system of inequalities [see [Figure 0.1](#)]. This proves that

$$x \cdot A + (1 - x) \cdot B, \quad \frac{1}{3} \leq x \leq \frac{1}{2}.$$

are undominated.

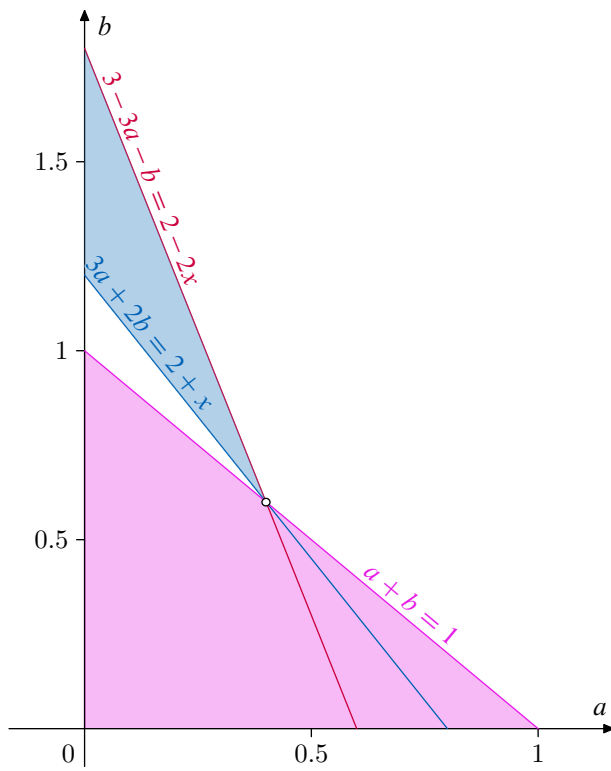


Figure 0.1: There is no solution ( $x = 0.4$ )